



Raft With RAPID

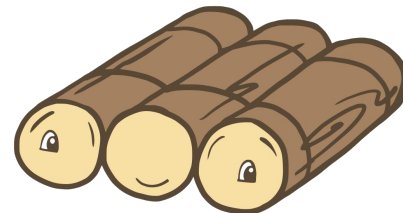
Creating a Simpler Consensus Algorithm

Presenters: Alex Chandler, David Klinger

12-04-2023

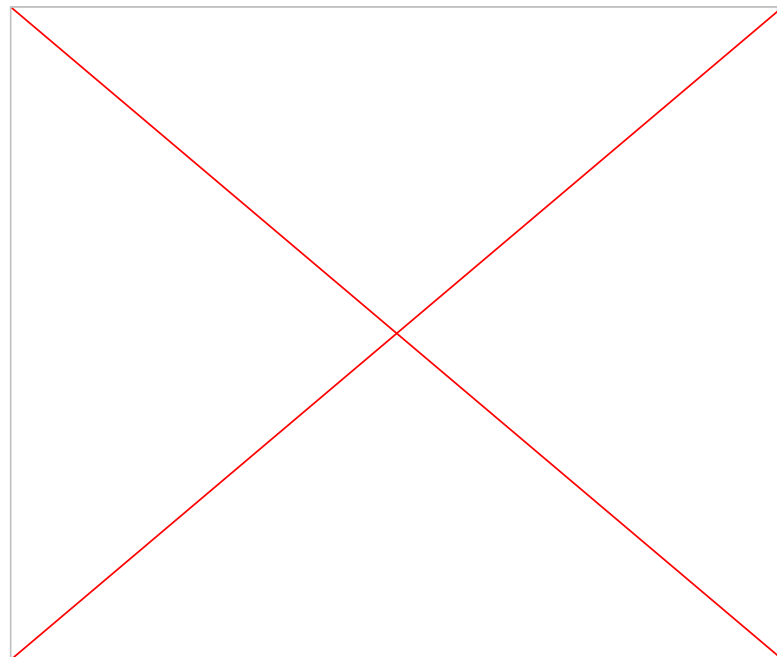
Motivation

- In Search of A Simpler Consensus Algorithm
- Why Build off of Raft?
 - Designed for Simplicity and Understandability:
 - Breaks down consensus into manageable subproblems.
 - More straightforward to learn, implement, and debug compared to complex algorithms like Paxos.
 - We want to make Raft Simpler
 - Widely Used with Strong Community Support:
 - Gained popularity for its clarity in handling consensus.
 - Extensive documentation, open-source implementations, and support across various programming languages.



Raft Overview

- Leader-Centric Model:
 - One leader manages log replication; followers accept entries.
 - Leader sends heartbeats to followers.
- Node States and Transitions:
 - Followers become Candidates if no heartbeat from leader.
- Leader Election with Randomized Timeouts:
 - Triggered when a follower's randomized timeout elapses.
 - Candidates request votes; wins leadership with majority.
- Terms and Consistency:
 - The cluster operates in terms of numbered “terms”.
 - Only one leader allowed per term to ensure consistency.
- Log Contents and Role:
 - Logs contain client commands and configuration changes.



RAPID Overview

- Why Use RAPID:
 - RAPID provides Stable and Consistent Membership Views.
- Multi-Process Cut Detection (CD):
 - Provides stable and accurate failure detection by requiring multiple confirmations for suspecting process faults.
- Leaderless Consensus Protocol:
 - Multi-Process Cut Detection ensures almost everywhere agreement.
 - Requires only one round of communication to decide the next configuration.
- Scalability and Speed:
 - Offers 2-5.8x faster cluster bootstrapping than Memberlist and ZooKeeper, while ensuring consistency and stability in large-scale deployments.
- Integration and Versatility:
 - Easily integrates into existing distributed applications.
 - Proven effectiveness in a range of enterprise data center environments and failure scenarios.



Stable and Consistent Membership at Scale with Rapid

Lalith Suresh, Dahlia Malkhi, and Parikshit Gopalan, *VMware Research*;
Ivan Porto Carreiro, *One Concern*; Zeeshan Lokhandwala, *VMware*

<https://www.usenix.org/conference/atc18/presentation/suresh>

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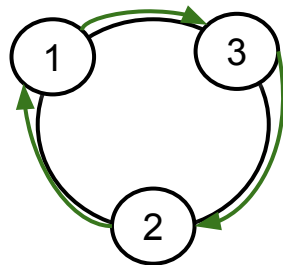
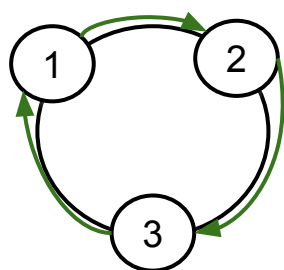
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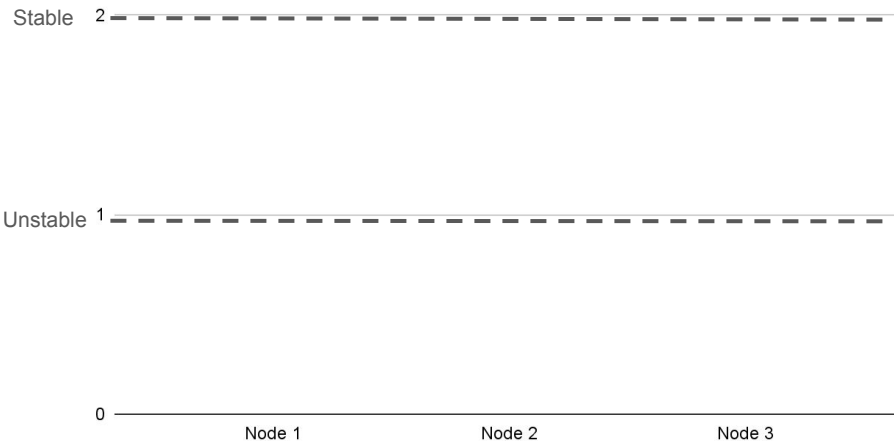


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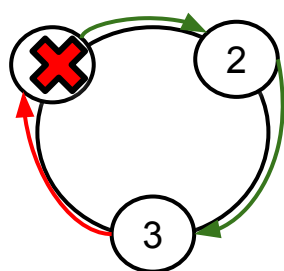
RAPID Fault Detection Mechanism



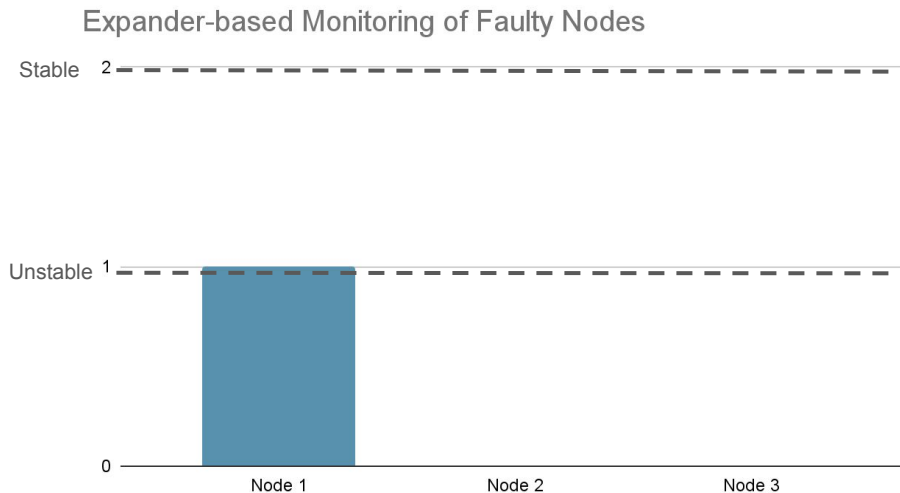
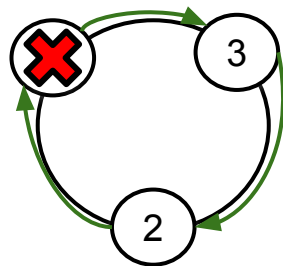
Expander-based Monitoring of Faulty Nodes



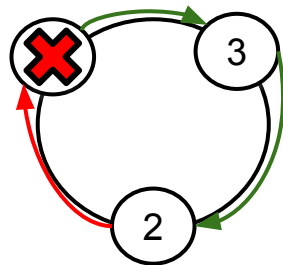
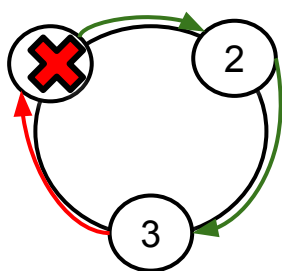
RAPID Fault Detection Mechanism



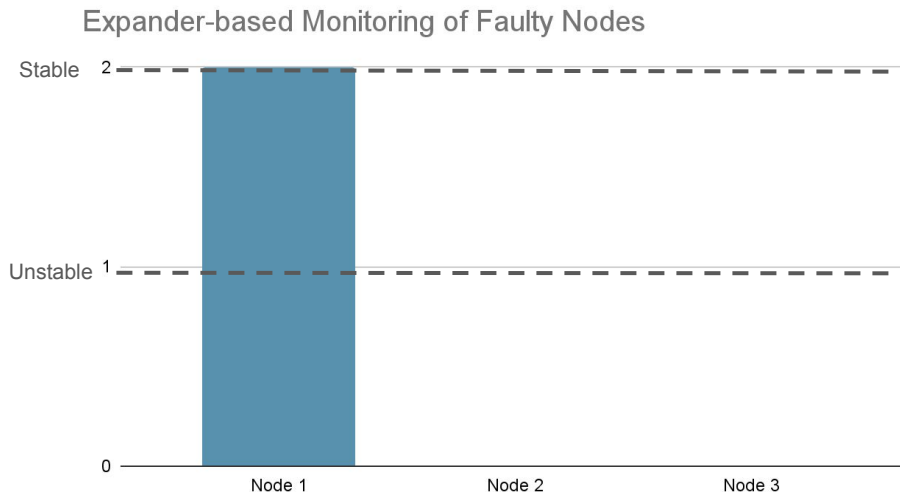
3 realizes first
that 1 is not
responding



RAPID Fault Detection Mechanism



2 then realizes
that 1 is not
responding



Advantages of RAPID in Raft: How RAPID simplifies and enhances Raft

- **Simplification of Cluster Management:**
 - RAPID assumes responsibility for tracking active nodes and managing cluster configurations, reducing the complexity of Raft's replicated log.
 - No longer need two overlapping cluster configurations before: (old→old_new→new) after: (old→new).
- **Simplifies Leader Implementation:**
 - Simplifies the leader's role by eliminating the need for continuous heartbeat messages, focusing only on proposing new log entries.
- **Simplifies Log Implementation:**
 - Raft stores both client values and cluster configuration changes in the log.
 - In Raft with RAPID, the replicated log now only stores client-proposed entries.

How Integrating RAPID Simplifies Cluster Configuration Changes

Before: Raft Configuration Change

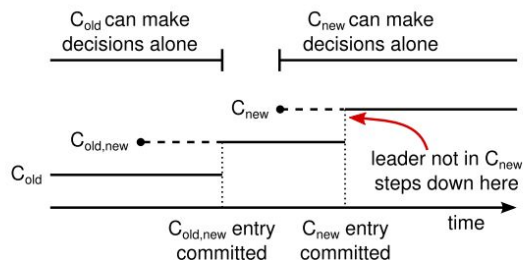
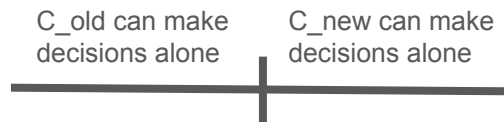


Figure 11: Timeline for a configuration change. Dashed lines show configuration entries that have been created but not committed, and solid lines show the latest committed configuration entry. The leader first creates the $C_{old,new}$ configuration entry in its log and commits it to $C_{old,new}$ (a majority of C_{old} and a majority of C_{new}). Then it creates the C_{new} entry and commits it to a majority of C_{new} . There is no point in time in which C_{old} and C_{new} can both make decisions independently.

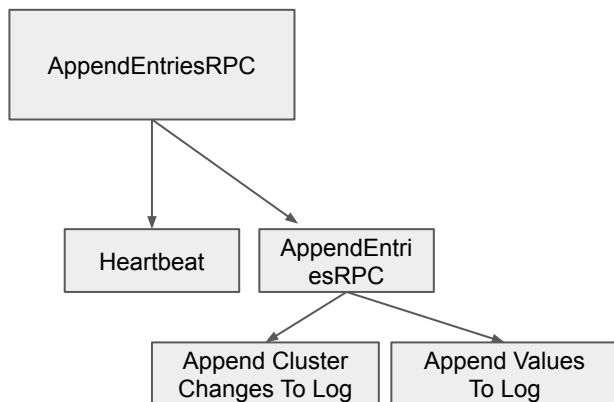
After: Raft-RAPID Configuration Change



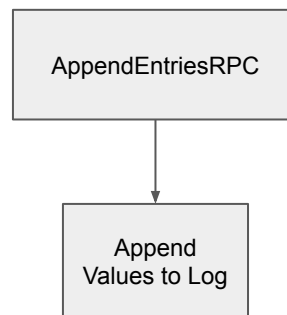
Limitation: The replication factor does not change

How Raft-RAPID Simplifies Handling of Append Entries RPC

Raft Append Entries



Raft-RAPID Append Entries

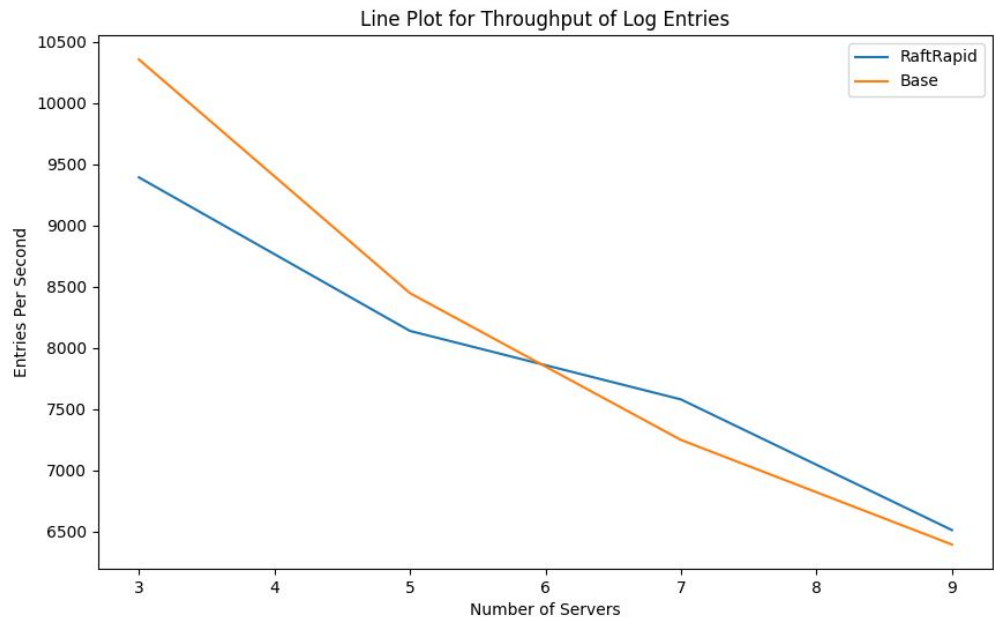


Evaluation

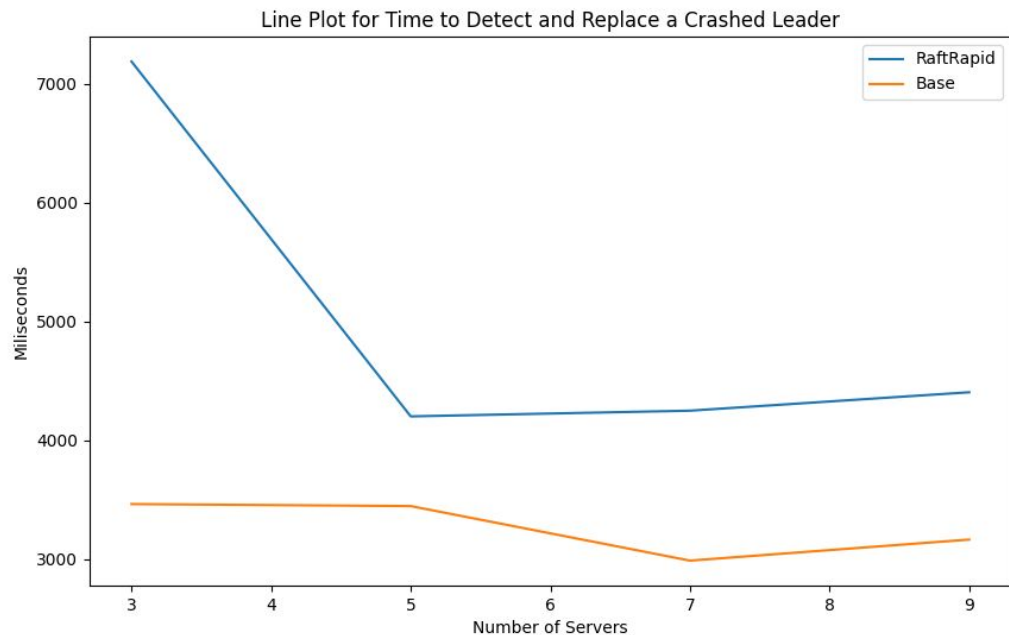
We compare Raft to Raft with RAPID along the following

1. Throughput of log entries
2. Time to detect and replace a crashed leader
3. Time from network partition healing to elected leader
4. Memory utilization
5. Network usage
 - a. Bandwidth
 - b. Latency
6. Configuration change time

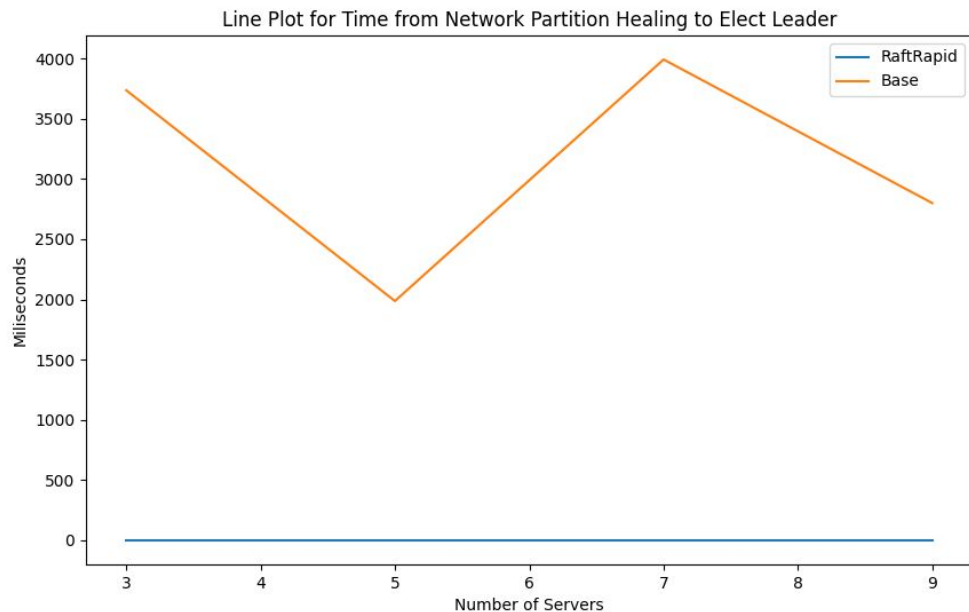
Throughput of Log Entries



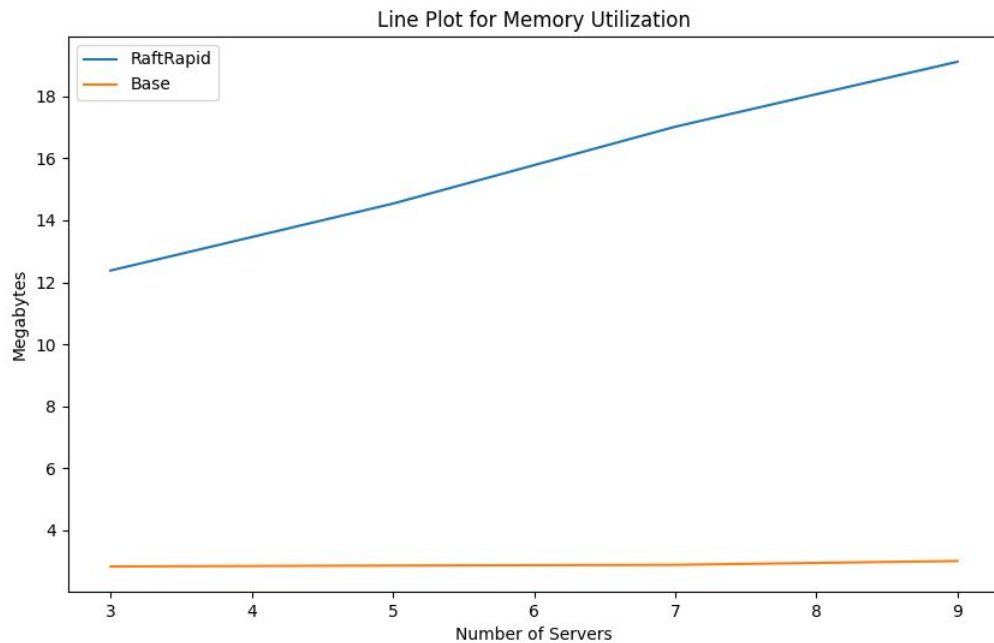
Time to Detect and Replace a Crashed Leader



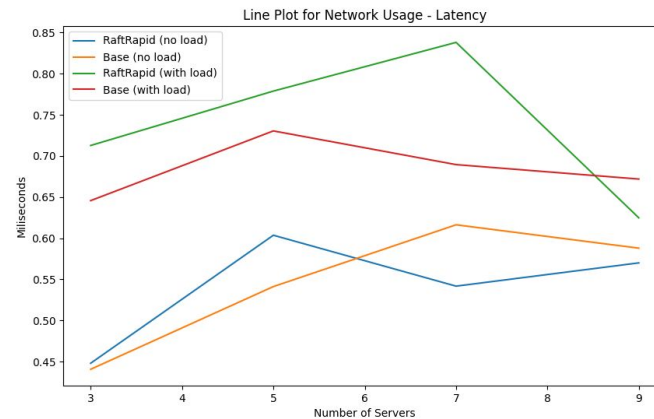
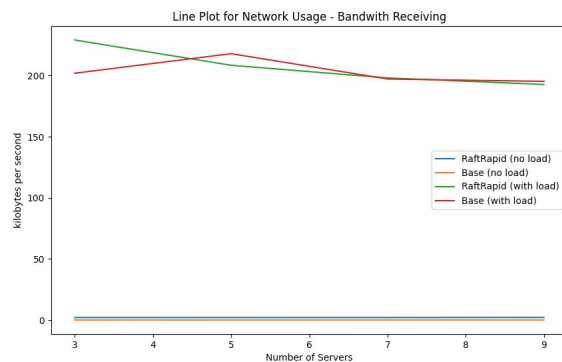
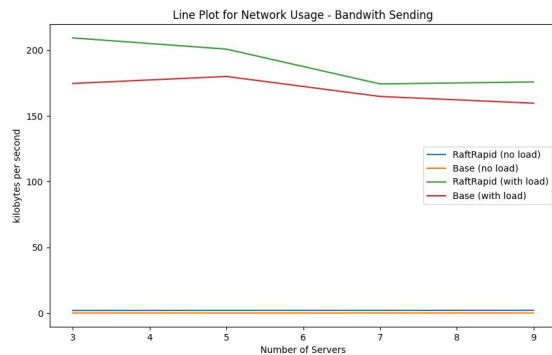
Time from Network Partition Healing to Elected Leader



Memory Utilization



Network Usage



Survey

Link in chat for voting Raft or Raft-RAPID:

Also available here: <https://forms.gle/1y2ZCaT1m3vydans5>

Thank You!